

## §19 Connecting with LS Theory

Lecture 11

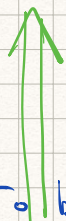
02/09

### - Critical Value selectors

Review minimax

This is the first application of  
of the min/max Principle

Goal: reprove



$$|\text{crit}(f)| \geq c_1(P) + 1$$

To be more precise we'll prove

$\text{crit}(f)$  is isolated

$$\Rightarrow cv(f) \geq c_1(P) + 1$$

• As before  
 $f: P \xrightarrow{c_0} \mathbb{R}$   
↑ closed

By a much more  
flexible construction

• Fix  $\mathbb{F} \leftarrow$  the background ring (or field)  
 $H_*(P) = H^{n-*}(P)$  when a field



## Critical value selectors

- $\alpha \in H_k(P) = H^{n-k}(P)$
- Stick to some construction of  $H_k(P)$  as the homology of  $C_*(P)$

$$\mathcal{F}_\alpha = \{A \mid [A] = \alpha\}$$

Remark •  $A$  is not necessarily a subset of  $P$  but rather a combination of mops to  $P$  ( $A: M \rightarrow P$  or simpl. chain) and then we mean the image of this map

- Or it can be liberally a subset:
  - pseudo-cycles [McDuff-Salamon]



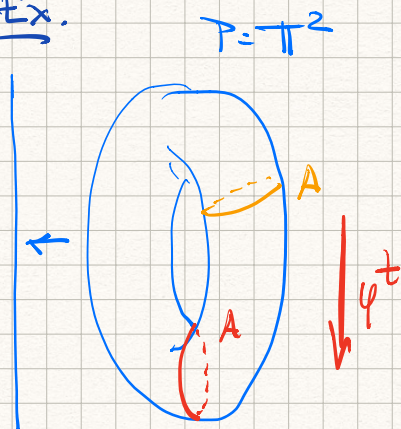
## Def. 1 Critical value selectors

$$c_\alpha(f) := C_{\mathcal{F}_\alpha}(f) := \min/\max \text{ for } \mathcal{F}_\alpha = \{A \mid [A] = \alpha\}$$

$$c_\alpha(f) = \inf_{[A] = \alpha} \max_A f$$



Ex.



$\alpha$  = homology class of these loops  
 $c_\alpha(f)$  is attained by pushing down  $A$  by  $\varphi^t \rightarrow 0$  till it "hangs on" a critical pt

Def 2 • Assume  $f$  is Morse ← a finite collection  
 $A = \sum m_i \alpha_i$ ,  $[A] = \alpha \in CM_\#(f)$

$$c_\alpha(f) = \min_A \left( \sum_{m_i \neq 0} \max f(x_i) \right)$$

- If  $f$  is not Morse, approximate it by Morse  $\tilde{f} \xrightarrow{C^0} f$  and set

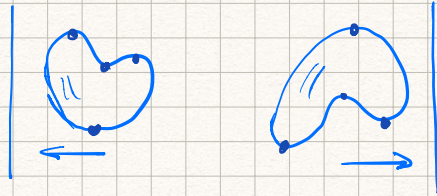
$$c_\alpha(f) = \lim_{\tilde{f} \rightarrow f} c_\alpha(\tilde{f})$$



Examples

- $c_{[M]}(f) = \max f$

- $c_{[pt]}(f) = \min f$



Def 3

$$c_\alpha(f) = \inf \{ a \mid \exists \text{ regular } \alpha \in \text{im}(H_* (P_a) \rightarrow H_* (P)) \}$$

Pf of equiv:

$$\text{Def 2} \Leftrightarrow \text{Def 3} \Leftrightarrow \text{Def 1}$$

by def                      by def  
when  $f$  is Morse

Properties

- Criticality:  $c_\alpha(f) = \text{crit value of } f$  (Min/Max)
- Monotonicity:  $f \leq g \Rightarrow c_\alpha(f) \leq c_\alpha(g)$  (By def)
- $C^0$  continuity:  $|c_\alpha(f) - c_\alpha(g)| \leq \underbrace{\|f - g\|_{C^0}}_{\max |f - g|}$

Hint: use  $f - \epsilon \leq g \leq f + \epsilon$   
and monotonicity

Remark  $c_\alpha(f) = f(x) \underset{\text{Crit}(f)}{\cap}$

cannot make  $x$  cont in  $f$



• sub-additivity  $\alpha, \beta \in M_x(\mathbb{P})$

$$(1) \quad c_{\alpha \cap \beta}(f+g) \leq c_\alpha(f) + c_\beta(g)$$

In particular  $\downarrow$  Take  $g=0$

$$(2) \quad c_{\alpha \cap \beta}(f) \leq c_\alpha(f)$$

In formal explanation: for (1)

$$\alpha = [A], \quad \beta = [B], \quad \alpha \cap \beta = [A \cap B]$$

← cycles →

$$c_{\alpha \cap \beta}(f+g) = \inf_{[C] = \alpha \cap \beta} \max(f+g)|_C$$

$$\begin{aligned} \text{take } C = A \cap B & \leq \max(f+g)|_{A \cap B} \\ & \leq \max f|_A + \max g|_B \end{aligned}$$

Now take inf over all  $A \& B$

• Can be turned into a pf once "cycles" and  $A \cap B$  are defined

(2) is even easier:

$$\max f|_{A \cap B} \leq \max f|_A$$



Pf of (i) using Y-graph flows

Can assume that  $f$  &  $g$  are Morse

$$\alpha \in \text{HM}_*(f) = H_*(P), \quad A \in \text{CM}_*(f)$$

$$\beta \in \text{HM}_*(g) = H_*(P), \quad B \in \text{CM}_*(g)$$

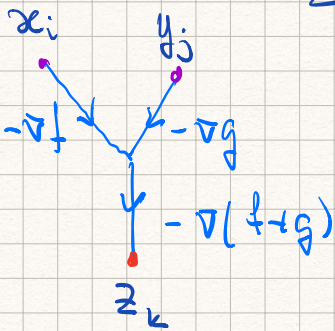
$$A = \sum_{\text{Crit}(f)} a_i x_i$$

$$B = \sum_{\text{Crit}(g)} b_j y_j$$

$$\alpha \cap \beta = [A \cdot B]$$

$$= \sum a_i b_j x_i \cdot y_j$$

$$= \sum a_i b_j \# \left\{ \begin{matrix} x_i & y_j \\ \swarrow & \searrow \\ & z_k \end{matrix} \right\} z_k$$



Pick  $A$  &  $B$  so that  
(exists)

$$c_\alpha(f) = \max f(x_i)$$

$$c_\beta(g) = \max g(y_j)$$

one of these  $C$

$(f+g)(z_k) \leq f(x_i) + g(y_j)$

Then  $c_{\alpha \cap \beta}(f+g) = \inf \max (f+g)$   
 $C = \sum c_k w_k$

$$\leq \max f(z_k)$$

$$\leq \max f(x_i) + \max g(y_j) = c_\alpha(f) + c_\beta(g)$$



Ex Give a direct pf of (2)  
using the  $M_x(P)$ -mod- str on  $MM_x(f)$

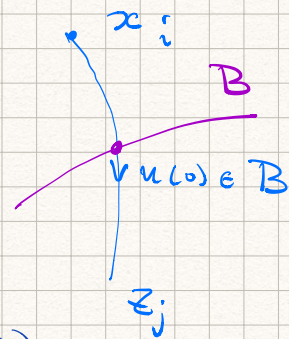
$$MM_x(f) \otimes M_x(P) \longrightarrow MM_x(f)$$

$$\alpha = [A] \quad \beta = [B]$$

$$A = \sum a_i x_i$$

$$B \cdot A = \sum a_i (B \cdot x_i)$$

$$= \sum a_i \# \left\{ \begin{array}{l} \text{ } \\ \text{ } \end{array} \right\} z_j$$



$f$  decreasing along  $u \Rightarrow (2)$

$\Delta$

---


$$\underline{\text{Cor}} \quad \min f \leq c_\alpha(f) \leq \max f$$



A more subtle result

Thm ("LS inequality")

Assume

- $\text{Crit}(f)$  isolated
- $|\beta| < n$  :  $\beta \in \mathbb{R}^n$

$\Rightarrow c_{\text{sup}}(f) < c_{\alpha}(f)$   
strict

I don't know a very simple pf

Explanation:

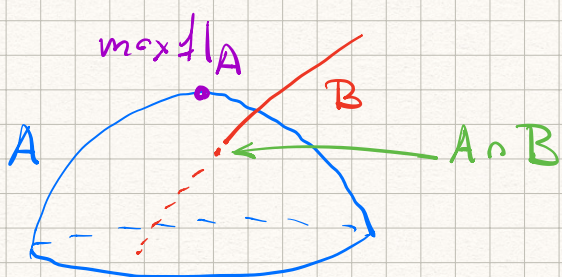
Assume  $c = \inf_{A} \max f|_A$  attained on some  $A$  (Not essential)

$\Rightarrow c = \max f|_A$  and  $\max f|_A$  is attained at one pt on  $A \Leftarrow \text{Crit}(f)$  isolated

$|\deg \beta| < n \Rightarrow \text{codim } B \geq 1$

$\Rightarrow B$  does not pass through  $\max$  of  $f|_A$  generically

$\Rightarrow \max f|_{A \cap B} < \max f|_A \triangleleft$



Not hard to turn into a real pt



Now we are ready to re-prove

Prop (LS)  $f: P \xrightarrow{C^2} \mathbb{R}$   
 $\text{crit}(f)$  isolated

$$\Rightarrow |cv(f)| \geq cl(P) + 1$$

Pf.  $k = cl(P)$ , use intersection product  $\cap$

$\beta_1 \cap \dots \cap \beta_k \neq 0$  in  $H_*(P)$   
 necessarily [pt] by PD  $|\beta_i| < n$   
not  $\sim$  units

$$\Rightarrow \underbrace{[M]}_{\alpha_0}, \underbrace{[M] \cap \beta_1}_{\alpha_1 = \alpha_0 \cap \beta_1}, \underbrace{[M] \cap \beta_1 \cap \beta_2}_{\alpha_2 = \alpha_1 \cap \beta_2}, \dots, \underbrace{[M] \cap \beta_1 \cap \dots \cap \beta_k}_{\alpha_k = \alpha_{k-1} \cap \beta_k}$$

Prop

$$\Rightarrow \underbrace{c_{\alpha_0}(f)}_{\text{max}} > c_{\alpha_1}(f) > \dots > \underbrace{c_{\alpha_k}(f)}_{\text{min}}$$

$k+1$  distinct values

$\triangleleft$



## Refinement:

Q. what happens when we have an equality of two crit value selectors?

$$c_{\alpha\beta}(f) \stackrel{\leq}{=} c_{\alpha}(f) \\ < \text{--- if isolated } |\beta| < n$$

Let  $K =$  crit set on the level  
 $f = c$ ,  $c = c_{\alpha\beta}(f) = c_{\alpha}(f)$   
 $\alpha, \beta \in H_*(P)$

Prop The restriction of  $PD(\beta) \in H^*(P)$   
(Ex) is  $\neq 0$  in  $H^*(K)$

Ex.  $\beta = [P] \Rightarrow PD(\beta) = 1 \in H^0(P)$

$$\begin{array}{ccc} H^0(P) & \xrightarrow{\text{onto}} & H^0(K) \\ 1 & \xrightarrow{\quad} & 1 \end{array} \quad \neq 0 \text{ always}$$

and  $\alpha\beta = \alpha$   
 $c_{\alpha\beta}(f) = c_{\alpha}(f)$

Cor  $|\beta| < n \Leftrightarrow |PD(\beta)| > 0$   
 $\Rightarrow \text{card}(K) = \text{continuum (cannot be countable)}$

Cor  $|cv(f)| \leq c(P)$   
 $\Rightarrow \text{continuum of crit pts}$



Rmk Need to be careful about the def  
of  $H^*(K)$ , for  $K$  could be a very bad  
set.



§ 20

Another Application:  
the Courant-Fischer Thm

Lecture 12  
02/11

setting

•  $Q(x) = \langle Ax, x \rangle : \mathbb{R}^{n+1} \rightarrow \mathbb{R}$ ,  $A^T = A$   
a quadratic form

•  $f = Q|_{S^n} : S^n \rightarrow \mathbb{R}$ .

Also known as the Rayleigh-Ritz quotient:

$$f(x) = \frac{\langle Ax, x \rangle}{\langle x, x \rangle} : \mathbb{R}^{n+1} \setminus \{0\} \rightarrow \mathbb{R}$$

Here we think of  $f$  as  $S^n \rightarrow \mathbb{R}$

•  $\lambda_0 \leq \dots \leq \lambda_n$  the eigenvalues of  $Q$  (or  $A$ )

Thm (Courant - Fischer)

$$\lambda_k = \inf_{\{L_k\}} \max_{L_k} f|_{L_k}$$

E.g.

$$\lambda_0 = \min f$$

$$\lambda_n = \max f$$

where  $L_k = S^n \cap \{ \text{dim lin space} \} \cong S^k$   
 $\uparrow$   $G_n(n+1, k+1)$

Rmk A similar statement for Hermitian forms on  $\mathbb{C}^{n+1}$

(Some  $p \neq 1$ )  
 $A = A^*$

E.g.  $k=1$   
 $\{L_0\} \cong \mathbb{R}P^1$



Pf.

- $\text{crit}(f) \leftrightarrow$  unit eigenvectors  
 $\text{cv}(f) \leftrightarrow$  eigenvalues

Pf

Lagrange multipliers

Crit of  $Q$  on  $\{g=c\} = \text{sol of } \nabla Q = \lambda \nabla g$

$$\nabla Q = 2Ax \quad (A^T = A)$$

$$\nabla g = 2x \quad g = \|x\|^2$$

$$\nabla Q = \lambda \nabla g : Ax = \lambda x, \quad \|x\| = 1$$

$\lambda$  is an eigenvalue

Then  $f(x) = Q(x) = \langle Ax, x \rangle = \lambda$

Remark

unit eigenvectors (=  $\text{crit}(f)$ )  
come in pairs  $\pm x$  with the  
same crit value  $\lambda$

- $\mathcal{F} = \{L_k\} \leftrightarrow G(n+1, k+1)$

To check: closed under  $\varphi^t$   
(requires a pf)

$\nabla Q$  is linear on  $\mathbb{R}^{n+1}$

$\rightarrow \nabla f = \text{proj of } \nabla Q|_{\mathcal{F}_k} \text{ to } \mathcal{F}_k$

$\rightarrow$  Need to check that  $\varphi^t$  sends  
equators to equators



min/max  $\Rightarrow \inf_{\mathcal{F}} \max f|_{L_k} = \mu_k$   
 are critical values

- Routine:  $\mu_k \leq \mu_{k+1}$  next page  
 Get  $\mu_0 \leq \dots \leq \mu_n : n+1$   
 $\lambda_0 \leq \dots \leq \lambda_n :$  Need to check that every c.v. is obtained in this way  
 $\mu_i$ 's are among  $\lambda_j$ 's  
 clear when distinct  $\mu_i = \lambda_j$  for a generic  $Q$   
 then pass to a limit  $\triangleleft$

Prmk. Rather than working on  $S^k$   
 could work on  $\mathbb{R}P^k : f(-x) = f(x)$   
 $\rightarrow$  Then  $L_k$  projects to  $\hat{L}_k \subset \mathbb{R}P^k$   
 linear proj subspace  $\cong \mathbb{R}P^k$

$$[\hat{L}_k] = [\mathbb{R}P^k] = \alpha_k \in H_k(\mathbb{R}P^k)$$

$\mathcal{F}_k \subset \mathcal{F}_{\alpha_k} \leftarrow$  as defined previously  
 all cycles in  $\alpha_k$

$$\lambda'_k = \inf_{\hat{L}_k \in \mathcal{F}_k} \max f|_{\hat{L}_k} \geq \inf_{A \in \mathcal{F}_{\alpha_k}} \max f|_A = \lambda_k$$

$\lambda'_0 \leq \dots \leq \lambda'_n$   
 $\lambda_0 \leq \dots \leq \lambda_n$  } both sequences formed by crit values (generically distinct)



Fixing the gap: all eigenvalues occur.

By continuity, assume  $\lambda_i$ 's are distinct

Assume that one of the eigenvalues does not occur

$$\begin{array}{ccc} \lambda_i < \lambda_{i+1} \\ \text{"} & & \text{"} \\ \mu_i = \mu_{i+1} \end{array}$$

Two ways to reason:

1) As in refinement of LS:

Then  $\text{Crit}(Q)$  are not isolated

2) By passing to  $\mathbb{R}P^n$

Then  $\{Q < \lambda_{i+1}\}$  is contractible  
to  $\mathbb{R}P^i$  and  
 $\{Q < \lambda_{i+1}\} \cong \mathbb{R}P^{i+1}$   
impossible  $\rightarrow \leftarrow$



§ 21 History Dimension:  
Lusternik-Schnirelmann Thm

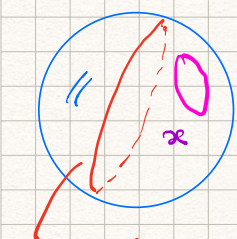
The origins of LS theory is

Thm (LS, 1929)

- Any metric on  $S^2$  has  $\geq 3$  simple closed geodesics
- Very difficult
  - Complete detailed pf: Ballmann 1978
- ↑  
embedded  
(no self-intersections)

Idea of the pf  $\leftrightarrow$  LS theory

- $\Lambda =$  smooth embedded loops in  $S^2$   
 (unoriented, unparametrized)



a great circle

$$S^1 \hookrightarrow S^2$$

For instance

$$\{\text{great circles}\} \subset \Lambda$$

$$\{S^2 \cap V\}$$

$$Gr(3,2) \cong \mathbb{RP}^2 \leftarrow \text{2-dim lin. subspace}$$

$$\mathbb{RP}^2 \xrightarrow{\sim} \Lambda \leftarrow \text{homotopy equivalent (absolutely mar. obs.)}$$



•  $L: \Lambda \rightarrow \mathbb{R}$  the length functional

By def  $\text{Crit}(L) =$  simple closed  
geodesics

• Punch line:  $c1(\Lambda) = c1(\mathbb{RP}^2) = 2$

If we could do LS on  $\Lambda$   
we would have

$$|\text{Crit}(L)| \geq c1(\Lambda) + 1 = 2 + 1 = 3$$

One can do an analogue of the LS  
theory on  $\Lambda$  (that was the starting pt  
for them), but this is very difficult.

◁



## §22 Getting rid of compactness: the Palais-Smale condition

Goal: Find a good replacement for the condition that  $P$  is compact  
Too restrictive.

Background assumption:

$P =$  finite dimensional (but possibly not compact)  
(preferable) on

Hilbert or Banach manifold  
(never compact)

$f: P \rightarrow \mathbb{R}$  sufficiently smooth

Def  $f$  satisfies the Palais-Smale condition (PS) if

• every seq  $\{x_i\} \subset P$  s.t.

$$\boxed{|f(x_i)| < c < \infty \quad \text{and} \quad df(x_i) \rightarrow 0} \\ (\Leftrightarrow \nabla f(x_i) \rightarrow 0)$$

contains a convergent subseq

Rmk  $df(x_i) \rightarrow 0 \Leftrightarrow \nabla f(x_i) \rightarrow 0$   
where  $\exists$  a R.m.:  $\dim P < \infty$  or  
 $P$  Hilbert  
but not when  $P$  is Banach.



Rmk PS condition does not involve the flow of  $-\nabla f$ . One can have PS satisfied without the flow

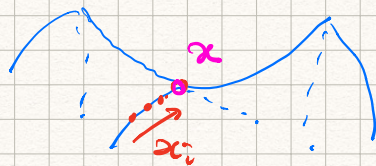
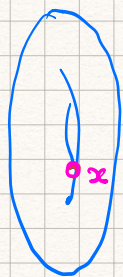
Important: PS condition easily breaks down

Don't need anything exotic

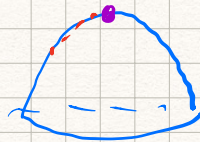
Construction: Start with  $f: \underset{\text{closed}}{P} \rightarrow \mathbb{R}$

•  $\dot{P} = P \setminus \{ \text{a crit pt} \}$

$\Rightarrow f: \dot{P} \rightarrow \mathbb{R}$  does not satisfy the PS but everything else is fine (  $\varphi_t$  is defined for all  $t$  )



$|f(x_i)|$  bounded  
 $\nabla f(x_i) \rightarrow 0$   
 but  $x_i$  does not converge  
 $x_i \rightarrow x \notin \dot{P}$



Upshot: Change of topology from  $P_0$  to  $P_1$   
Without PS  $\Rightarrow$  ~~existence of critical pts~~



## Basic example - application - illustration

### Assumptions:

- $f$  bounded from below
  - $f$  satisfies PS
  - The flow for anti-grad-like vector field is defined for all  $t \geq 0$
- $\Rightarrow f$  has a critical pt actually min is attained

$$X \text{ s.t. } L_X f \leq 0$$

$$L_X f(x) = 0 \Leftrightarrow x \in \text{Crit}(f)$$

E.g.  $X = -\nabla f$ .

For many purposes as good as  $-\nabla f$  but more flexible & does not need metric

Pf • Start with some  $x$  and set

$$z_k = \varphi^k(x_0) \quad k \rightarrow \infty$$

PS + bounded from below

$$\Rightarrow z_k \rightarrow y \leftarrow \text{crit pt}$$

- To get min: take  $x_i$  s.t.

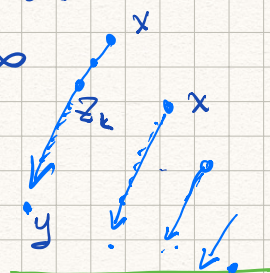
$f(x_i) \rightarrow \inf f$ . Apply this process and get  $y_i$

$$\text{As } df(y_i) = 0 \text{ \& } f(y_i) \rightarrow \inf f$$

$$\Rightarrow y_i \rightarrow y \text{ \& } f(y) = \inf f$$

$\triangle$

(135)





## Back to Minimax

### Setting

- $P$  as above
  - $f: P \rightarrow \mathbb{R}$  satisfies PS
  - the anti-grad flow  $\varphi^t$  of  $f$  is defined for all times  $t \geq 0$
  - $\mathcal{A}$  is closed under  $\varphi^t$
  - $f$  on  $\mathcal{A}$  is bounded from below:  
$$\inf_{A \in \mathcal{A}} \max f|_A > -\infty$$
- With some care can replace by anti-grad-like

Thm (Minimax, II)  $c = \inf_{A \in \mathcal{A}} \max f|_A$  is a crit value

Ex. Take  $\mathcal{A} = \text{pts of } P$  in the previous example.

The pt is very similar to the example



Pf • Bounded from below  $\Rightarrow c > -\infty$

$\in \mathbb{R}$

• Will prove:  $\forall \varepsilon > 0 \exists \delta > 0 \exists x \text{ s.t. } c - \varepsilon < f(x) < c + \varepsilon \ \& \ \|\Delta f(x)\| > \delta$

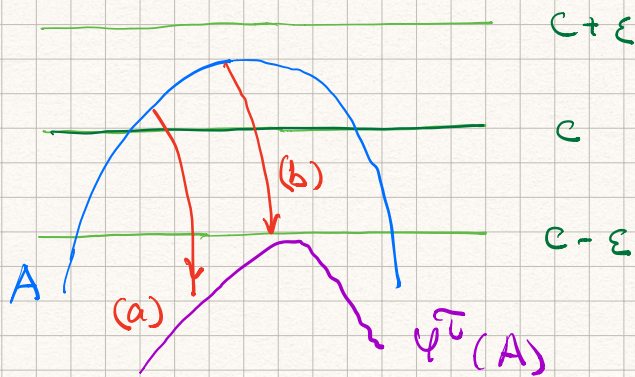
• Arguing by contradiction, assume that:

$$\|\Delta f(x)\| > \varepsilon \text{ when } c - \varepsilon < f(x) < c + \varepsilon$$

for some  $\varepsilon > 0$

•  $\exists A \in \mathcal{A}$  s.t.

$$\max_{f|_A} > c + \varepsilon$$



• claim  $\tau = \frac{2}{\varepsilon} \Rightarrow \max_{f|_{\psi^\tau(A)}} < c - \varepsilon$

Indeed, take  $x \in A$

(a) if  $f(\psi^t(x)) < c - \varepsilon$  for some  $t \in [0, \tau)$  we are done for  $f|_{\psi^\tau(A)}$



(b) if  $c - \varepsilon < f(\psi^t(x)) < c + \varepsilon$   
 for  $t \in [0, \varepsilon]$ , we have

$$|\nabla f(\psi^t(x))| > \varepsilon$$

$$\Rightarrow f(\psi^t(x)) < \underbrace{c + \varepsilon - \varepsilon^2 t}$$

$$\frac{d}{dt} f(\psi^t(x)) \leq -\langle \nabla f, \nabla \psi \rangle \leq -\varepsilon^2$$

$$\Rightarrow f(\psi^\varepsilon(x)) < c + \varepsilon - 2\varepsilon < c - \varepsilon$$

$$\Rightarrow \max_{\psi^t(x)} f < c - \varepsilon$$

△